

# **DK-MICROAGGREGATION: ANONYMIZING GRAPHS WITH DIFFERENTIAL PRIVACY GUARANTEES**

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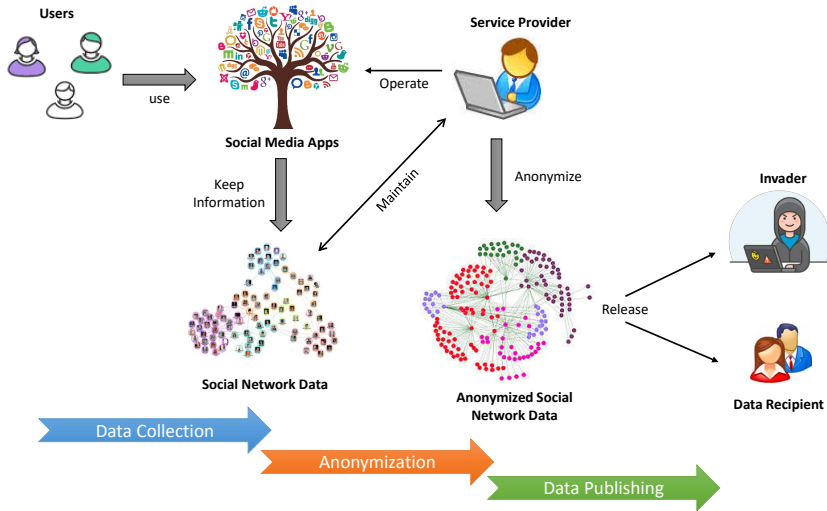
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# INTRODUCTION

- Graph data analysis has been widely performed in **real-life applications**. For instance,
  - ▶ **online social networks** are explored to analyze human social relationships;
  - ▶ **election networks** are studied to discover different opinions in a community.
- However, such networks often contain **sensitive** or personally identifiable information, such as **social contacts**, **personal opinions** and **private communication records**.
- Publishing graph data can thus pose a *privacy threat*.

# GRAPH DATA RELEASE PROCESS



**Figure 1:** Graph Data Release Process (e.g. online social network)

- **Aim:** To generate **anonymized graphs** with  **$\epsilon$ -differential privacy** guarantee for improving utility of anonymized graphs being published.
- **Key Challenges:**
  - ▶ To preserve **topological structures** of an original graph at different levels of granularity.
  - ▶ To enhance **utility** of graph data by reducing the magnitude of noise needed to achieve  $\epsilon$ -differential privacy through adding controlled perturbation to its edges (i.e., **edge privacy**).
- **Key Observation:** We observe that the **dK-graph model** [5] for analyzing network topologies can serve as a good basis for generating **structure-preserving** anonymized graphs.

# PROBLEM FORMULATION

- The **dK-graph model** [5] provides a systematic way of extracting subgraph degree distributions from a given graph, i.e. **dK-distributions**.

### DK-DISTRIBUTION

A *dK-distribution*  $dK(G)$  over a graph  $G$  is the probability distribution on the connected subgraphs of size  $d$  in  $G$ .

- Specifically, **1K-distribution** captures a degree distribution, and **2K-distribution** captures a joint degree distribution. When  $d = |V|$ , dK-distribution specifies the entire graph.
- A **dK-distribution** is extracted from a graph, by using **dK function** (s.t.  $\gamma^{dK}(G) = dK(G)$ ).



- We define *dk-graph* as a graph that can be constructed through reproducing the corresponding *dk-distribution*.

### *DK-GRAPH*

A *dk-graph* over  $dk(G)$  is a graph in which connected subgraphs of size  $d$  satisfy the probability distribution in  $dk(G)$ .

- Conceptually, a *dk-graph* is considered as an *anonymized version* of an original graph  $G$  that retains certain topological properties of  $G$  at a chosen level of granularity.
- We aim to generate *dk-graphs* with  $\epsilon$ -*differential privacy* guarantee for preserving privacy of structural information between nodes of a graph (*edge privacy*).

- Two graphs  $G = (V, E)$  and  $G' = (V', E')$  are said to be **neighboring graphs**, denoted as  $G \sim G'$ , iff  $V = V'$ ,  $E \subset E'$  and  $|E| + 1 = |E'|$ .

## Differentially private dK-graphs

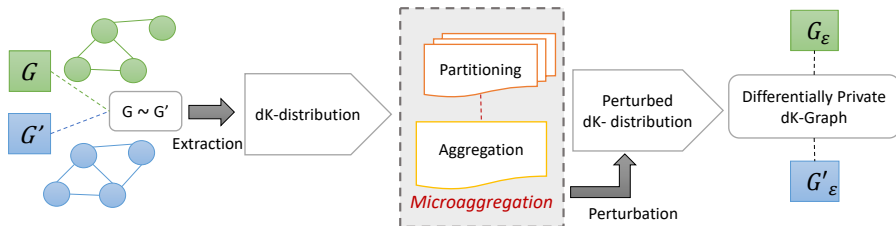
A randomized mechanism  $\mathcal{K}$  provides  $\varepsilon$ -differentially private dK-graphs, if for each pair of neighboring graphs  $G \sim G'$  and all possible outputs  $\mathcal{G} \subseteq \text{range}(\mathcal{K})$ , the following holds

$$\Pr[\mathcal{K}(G) \in \mathcal{G}] \leq e^\varepsilon \times \Pr[\mathcal{K}(G') \in \mathcal{G}]. \quad (1)$$

- $\mathcal{G}$  is a family of **dK-graphs**, and  $\varepsilon > 0$  is the **differential privacy parameter**. Smaller values of  $\varepsilon$  provide stronger privacy guarantees.

# **DK-MICROAGGREGATION FRAMEWORK**

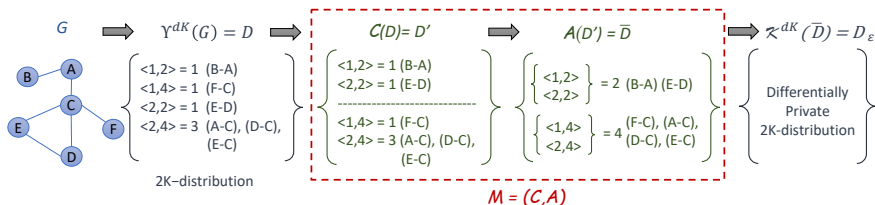
- We incorporate **microaggregation** techniques [1] into the **dK-graph model** [5] to reduce the amount of **random noise** without compromising  $\epsilon$ -differential privacy.
- Generally, **dK-microaggregation** works in the following steps:
  - (1) **Extracts** a dK-distribution from each *neighboring* graph.
  - (2) **Microaggregates** the dK-distribution and perturbs the microaggregated dK-distribution to generate  $\epsilon$ -differentially private dK-distribution.
  - (3) **Generates**  $\epsilon$ -differentially private dK-graphs using a dK-graph generator [4, 5].



**Figure 2:** A high-level overview of the proposed framework (dK-Microaggregation).

# PROPOSED ALGORITHM

- A **microaggregation algorithm** for dK-distributions  $\mathcal{M} = (\mathcal{C}, \mathcal{A})$  consists of two phases:
  - (a) **Partition** - similar tuples in a dK-distribution are partitioned into the same cluster;
  - (b) **Aggregation** - the frequency values of tuples in the same cluster are aggregated.



**Figure 3:** An illustration of our proposed algorithms.

- **MDAV-dK algorithm:** We use a simple microaggregation heuristic, called *Maximum Distance to Average Vector (MDAV)* [1], which can generate clusters of the **same size**  $k$ , except one cluster of size between  $k$  and  $2k - 1$ . Then unlike MDAV, we aggregate frequency values of tuples in each cluster.

However, **MDAV-dK** would suffer significant information loss when evenly partitioning a **highly skewed**  $dK$ - distribution into clusters of the same size.

- **MPDC-dK algorithm:** To address this issue, we propose *Maximum Pairwise Distance Constraint (MPDC-dK)*, which aims to partition a  $dK$ -distribution into a minimum number of clusters in which every pair of tuples from the same cluster satisfies a **distance constraint**  $\tau$ .



# EXPERIMENTS AND RESULTS

## ■ Three network datasets:

- (1) *polbooks* contains 105 nodes and 441 edges.
- (2) *ca-GrQc* contains 5,242 nodes and 14,496 edges.
- (3) *ca-HepTh* contains 9,877 nodes and 25,998 edges.

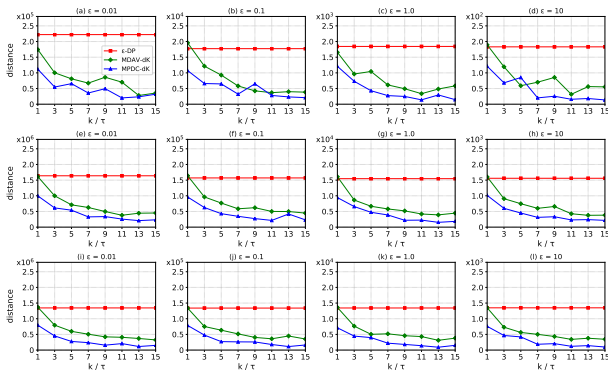
## ■ Two measures:

- ▶ **Euclidean distance** [6] measures network structural error between original and perturbed dK-distributions.
- ▶ **sum of absolute error** [2] measures within-cluster homogeneity of clustering algorithms, defined as:

$$SAE = \sum_{i=1}^N \sum_{\forall x_j \in c_i} |x_j - \mu_i|,$$

where  $c_i$  is the set of tuples in cluster  $i$  and  $\mu_i$  is the mean of cluster  $i$ .

- To verify the **utility**, we compare the structural error between original and perturbed dK-distributions generated by MDAV-dK, MPDC-dK and the baseline method  $\epsilon$ -DP. Our proposed algorithms **MDAV-dK** and **MPDC-dK** lead to **less structural error** for every value of  $\epsilon$  as compared to  $\epsilon$ -DP.



- We compare the **quality of clusters**, in terms of within-cluster homogeneity, generated by MDAV-dK and MPDC-dK. **MPDC-dK outperforms MDAV-dK** by producing clusters with less SAE over all three datasets.

**Table 1.** Performance of MDAV-dK under different values of  $k$ .

Datasets	Measures	$k=1$	$k=3$	$k=5$	$k=7$	$k=9$	$k=11$	$k=13$	$k=15$
<i>polbooks</i>	SAE	0	144.6	184.67	<b>224.84</b>	273.6	292.21	299.15	334.25
	# Clusters	161	53	32	<b>23</b>	17	14	12	10
<i>ca-GrQc</i>	SAE	0	1073.3	1476	<b>1810.5</b>	2166.8	2313.7	2555.5	2730
	# Clusters	1233	411	246	<b>176</b>	137	112	94	82
<i>ca-HepTh</i>	SAE	0	968.72	1304	1599.8	<b>1893.9</b>	2063	2232.9	2389.7
	# Clusters	1295	431	259	185	<b>143</b>	117	99	86

**Table 2.** Performance of MPDC-dK under different values of  $\tau$ .


Datasets	Measures	$\tau=1$	$\tau=3$	$\tau=5$	$\tau=7$	$\tau=9$	$\tau=11$	$\tau=13$	$\tau=15$
<i>polbooks</i>	SAE	90.72	<b>192.15</b>	328.96	424.2	563.73	617.63	723.06	795.77
	# Clusters	68	<b>25</b>	13	8	7	5	3	3
<i>ca-GrQc</i>	SAE	725.38	<b>1732.1</b>	2630.6	3470.6	4262.9	5176.7	6170.1	7037.7
	# Clusters	483	<b>178</b>	98	61	42	35	26	20
<i>ca-HepTh</i>	SAE	841.87	<b>1761.8</b>	2773.3	3721.4	4719.2	5623.8	6402.6	7034.2
	# Clusters	412	<b>140</b>	73	37	34	24	19	15

# CONCLUSION AND FUTURE WORK

## ■ Conclusion:

- ▶ We present a novel framework, called *dK-microaggregation*, that can leverage a series of network topology properties to generate  $\epsilon$ -differentially private anonymized graphs.
- ▶ We propose a **distance constrained algorithm** for approximating dK-distributions of a graph via microaggregation within the proposed framework, which can reduce the amount of noise being added into  $\epsilon$ -differentially private anonymized graphs.
- ▶ The effectiveness of our proposed framework has been empirically verified over **three real-world network**.

- **Future work:** To this work will consider **zero knowledge privacy (ZKP)** [3], to release statistics about social groups in a network while protecting privacy of individuals.

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**THANKS FOR YOUR ATTENTION!**

**ANY QUESTIONS**

